V, Queenhun Error Correction

1. lubroduction

a) Setting & Proten

· Complety to environment induces errors (i.e., uncontrolled Schandr). · Classical computes: Information stored in "un croscopie" properties - errors un likely. · Quantum computers: - need grøts = " sneple" grænhen systems, and und store general superposition, not just 107 and 117 -> fragile! - should be well isolated to protect qubits, but also reced compley to "exurament" (experimental apparetus) to control the computation (gates, manuales).

Q' Can we product quantum information from norse?

Clashical error correction:

Copy monualin, e.g. encode 1 bit m 3 bits: "encoding" 1 - 1:= 111 Error model: Bost plip 4/ some (small) probability P (mælependenky on all bik): => hypically O or 1 boks flipped. Error correction ("decodary") by majority vok: 000,001,010,100 -> 000 MA, MO, NO1, OH >> MA Probebility for a Cogstal ever "(i.e. on encoded bit): Peror = prob( = 2 flips) = p<sup>3</sup> + 3 p<sup>2</sup>(1-p)  $= 3p^2 - 2p^3 < p$  for p < |2.error quadratically suppressed!

= efective ever probability decreased.

Can be support by: - using more tob: 01 00...0, 11 11...1 - ueshy (" concertue hing") codes - using matter codes (i.e. ucode several bib at once)

Que hun Error Correctia:

Several potential prostems shew trying to flueralise classical error correction codes:

· caused copy gubib · even if we could : what would be the " majonly work"?

o different types of errors east,

e.g. X (61 flip) or 2 ("plax /lip")

· errors can be continuens: there is an

infruity of corors !

· measury qubits destroys quanter information."

6) The 3-qubit bit flip code

Copy qubits in computational basis:  $|o\rangle \leftrightarrow |o\rangle = |ooo\rangle$  $|1\rangle \longleftrightarrow |\hat{\lambda}\rangle = |11\rangle$ rie. Hu cucodites is a linear mos 14)= 2/0)+ B/1) - 2/000) + B/111> Possike encoding circuit:  $\frac{|\psi\rangle}{|\psi\rangle} = \frac{|\psi\rangle}{|\psi\rangle} + \frac{||||\rangle}{|\psi\rangle} = \frac{|\psi\rangle}{|\psi\rangle} + \frac{|||||\rangle}{|\psi\rangle} + \frac{|||||}{|\psi\rangle} + \frac{|\psi\rangle}{|\psi\rangle} + \frac{|\psi\rangle}{$ Nos consider but flip error on quot i: It's Kilt's Xilt's

Can we correct for one but pip error on an unknown quat i?

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Proklen; Meaning the gutits in comp. basis reveals i, but also destroys oupspoots?

- Need a measurement which only chines upor ache about posshow i of error - indep. of encoded state 14>!

Define "gudrone measurement" with outcomes 0,1,2,3, and projectors:

0 = "uo flip": Po = 1000×0001 + 1111 ×111 1 = "let gubit plipped": P1 = 1,000 × 100 1 + 1011 × 011 1  $P_2 = 1010 \times 010 + (101 \times 100)$ 2= 2nd gubt Kipped": P3 = 1001 X001/ + 1100 X1101 3 = 'Srd quest flipped :

(Ruis depues a complete reconcurrent, as ZP-=I)

The onterene is called the "error syndrome",

Resput of EPal reveals only 25th of rups, = one qu'est of monaha untouched!

268 By direct repector: The representa altered is the lacka of the bit flip, e.g. alooo) + \$ (111) = alooo) + \$ (101) on gubt 2 = D measurement always chins P2, with post - uncasure unent stake ~ (010) + \$ / lo1) (covery: ~ ~ (000) + \$ (111) . Kip gutt 2 = P Bot flip corrected! Works for any mule bot Kip muchan location and no plop, and for all states 147 = o supression of error p~> 3p<sup>2</sup>-2p<sup>3</sup>, as classically. By literaity, Kus also works for part of a lager enten fled that: x/0>/a>+ \$(1>/6> = x 1000> (a>+ \$/101>/6> x, x /100> (a>+ 8 /011>/6> meas.: P₁ → x/000>/a>+ Correct: X, SIAM >13>

What about conditions errors, e.g.  

$$|\psi\rangle \mapsto e^{i\partial X_i} |\psi\rangle = (\cos \partial I + i \sin \partial X_i) |\psi\rangle$$
?  
 $|\psi\rangle = \alpha(000) + \beta(100) \xrightarrow{error} \alpha(\cos \partial (000) + i \sin \partial (001))$   
 $e.g.X_3 \mapsto \beta(\cos \partial I + i \sin \partial (001))$ 

$$e.g. X_{3} + \beta (cod | M1 > + i m d (110))$$

$$= cod (a(000) + \beta | M1 >) + i sn d (a(00) + \beta | M0))$$

$$\int syndrme P_{0} \int syndrme P_{3}$$

$$prob.(and)^{2} prob.(18n \theta)^{2}$$

$$p = co^{2} \vartheta$$
: result  $P_{0}$ ,  
 $post-meas$ . Aak  $\alpha | 0000 \rangle + \beta | || 1 \rangle$ ,  
 $0 \equiv mo correcta$ :  
 $O \equiv Mo correcta$ :

$$p = 8\pi^2 \vartheta : rsult P_3,$$

$$port - meas. Date  $\ge 1001 > + \beta / 100 >,$ 

$$3 = correction : flip bit 3 :$$

$$\implies \alpha | 000 > + \beta | 111 > : (OKI)$$$$

Respersivent of error syndrome EPa's collapses 4 correctable continuous error onto one of the discrete errors: · maturement "dipitelires" error . reficient to study discrete (destrupishable) errors ( ush be formalized later) A different perspective on syndrome uncasurement & concha ( He "Mabilion porualion" - more lats): +1 eigenstates of 2722 and 223 1000/11/17: ("stabilizers") Reapure Zitz and Zz Zz: Campare qubits 182 and 283 (+1, +1): us ever (-1,+1): qubit 1 Kipped (H, -1) : gubt 3 Kapped (-1,-1) : gubit 2 Kipped

Row formally: eucoded stake  $|\psi\rangle = d(000) + \beta(100):$   $\implies 2_1 2_2 (\psi) = |\psi\rangle, 2_2 2_3 |\psi\rangle = |\psi\rangle$ Bit  $|\psi\rangle$  error, e.g.  $X_1$ :  $X_1$  sub-commute  $\ln K_1 = 2_1 2_2$   $\implies \langle \psi | X_1 + 1 + 2_1 X_1 | \psi\rangle = - \langle \psi | + 2_1 2_2 | \psi\rangle$ = -1

Rus: Outcome -1 for ZiZz (=) an error which anti-commutes work ZiZz has occurred,

The correction operation unit satisfy the same auh - Couran takt richans (and some futur proposhes) -> (aks!

Have focused on Xerros. Port chat about 2 cros? ~ (000) + \$ /111) ~ 2000) - \$ /111) ~ 2000) - \$ /111) Ruis is spill a state in the code space (i.e., a valid encoded state (12) => error unt dekchalle, but it has changed 14). After decoding, the ever ach as a/0> +\$(1) ~~> a/0>-\$11>, r.e. as a lopical 2 operation, "lopical operation" = oproha a encoded gubit. = 3-quest but plip code caund protect agailist siefle "please plip error" Z'.

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273 Stabilize picture: Error Zi communtes with stabilities 272 8722 = it cannot be dekckd. but: 2: cannot le capassed as a product of the stabilitors Etz l Eztz = D lopial error! c) The 3-qubit phase-flip coole Can we correct against 2 errors? チ |+> = 1-7, モ |->= |+> => Zerror = Sit fly error ~ |±>-5007. Use encody alos+\$117 -> alôs+\$12, unk los:= (+++>, lis= (--->: Will protect against surfle 2 errors! Eucodiz: 147 - H-107 - H-10) - 田-

Syndrome measurement:

 $\widetilde{P}_{a} := H^{\otimes 3} P_{a} H^{\otimes 3}$ 

(or vie debiliters X, K2 (X2X3).

Kecovery operation:

 $HX_i H = Z_i$ (anti- con. with stabilities).

Problem; Now, Kere is no protection against but thip errors X' - and Xi acts as a logical X operator