## Lecture & Proseminar 250120/250122 "Quantum Information, Quantum Computation, and Quantum Algorithms" WS 2020/21

— Exercise Sheet #10 —

## Problem 1: Phase estimation

Consider a unitary U with an eigenvector  $U|\phi\rangle = e^{2\pi i\phi}|\phi\rangle$ . Assume that

$$\phi = 0.\phi_1\phi_2...\phi_n = \frac{1}{2}\phi_1 + \frac{1}{4}\phi_2 + ... + \frac{1}{2^n}\phi_n$$

i.e.  $\phi$  can be exactly specified with *n* binary digits. Our goal will be to study ways to determine  $\phi$  as accurately as possible, given that we can implement U (and are given the state  $|\phi\rangle$ ).

- 1. First, consider that we use controlled-U operations  $CU|0\rangle|\phi\rangle = |0\rangle|\phi\rangle$ ,  $CU|1\rangle|\phi\rangle = |1\rangle e^{2\pi i\phi}|\phi\rangle$ . Describe a protocol where we apply CU to  $|+\rangle|\phi\rangle$ , followed by a measurement in the  $|\pm\rangle$  basis, to infer information about  $\phi$ . Which information, and to which accuracy, can we obtain with N iterations? (Bonus question: Could this scheme be refined by changing the measurement?)
- 2. Now consider a refined scheme. To this end, assume we can also apply controlled- $U^{(2^k)} \equiv CU_k$  operations for integer k efficiently.

a) We start by applying  $CU_{n-1}$  to  $|+\rangle|\phi\rangle$ . Which information can we infer? What measurement do we have to make?

b) In the next step, we apply  $CU_{n-2}$ , knowing the result of step a). What information can we infer? What measurement do we have to make? Rephrase the measurement as a unitary rotation followed by a measurement in the  $|\pm\rangle$  basis.

c) Iterating the preceding steps, describe a procedure (circuit) to obtain  $|\phi\rangle$  exactly. How many times do we have to evaluate controlled- $U^{(2^k)}$ 's?

(Note: This procedure is known as quantum phase estimation.)

## Problem 2: Factoring 15

Verify the factoring algorithm (i.e., the reduction to period finding described in the lecture – subsction 3.c) for N = 15 – i.e., consider all a = 2, ..., N-1, check wether gcd(a, N) = 1, find r s.th.  $a^r \mod N = 1$  (you don't have to use a quantum computer), and check if this can be used to compute a non-trivial factor of N. How many different cases do you find? What possible periods r appear?

## Problem 3: Controlled gates and measurements.

An *n*-qubit Toffoli gate is a Toffoli gate with n - 1 controls; i.e., it flips the *n*'th bit if and only if the other n - 1 bits are all one.

- 1. Show that the *n*-qubit Toffoli gate can be implemented using two n 1-qubit Toffoli gates and two regular 3-qubit Toffoli gates using one ancillary qubit.
- 2. Decomposing every gate into 3-qubit Toffoli gates, how many 3-qubit Toffoli gates do you need to construct the *n*-qubit Toffoli gate?
- 3. Find a construction which is more efficient in terms of the scaling of the number 3-qubit Toffoli gates used, at the cost of using more ancillas. (A linear number of 3-fold Toffoli gates should suffice.)

(*Hint:* Remember that the Toffoli gate can be used to build a logical AND gate using ancillas.)