

3. Applications of entanglement:

Teleportation and dense coding

a) Teleportation

Schup:



- A & B share entangled state $|\phi^+\rangle_{AB} = \frac{1}{\sqrt{2}}(|00\rangle_{AB} + |11\rangle_{AB})$
- A has unknown quantum state

$$|X\rangle_{A'} = a|0\rangle_{A'} + b|1\rangle_{A'}$$

(Could e.g. also be part of a larger system $\rightarrow$ security!)

- A & B cannot (reliably)* transmit quantum states, but can communicate classically "for free",

* If the line is unreliable, A & B can still use it to create entangled states $|\phi^+\rangle$, e.g. by repeat-until-success, or entanglement distribution ($\rightarrow$ later!), or using "quantum repeaters" ($\rightarrow$ later!)

Question: Can A get $|X\rangle$ (safely) to B?

Problem: Any measurement of $|K\rangle$ would ^{Chapter III, pg 21} only reveal partial information, yet destroy state!

Solution: Quantum Teleportation!

Teleportation Protocol:

① A performs measurement on $A'A$ in Bell basis

$$|\phi^+\rangle = \frac{1}{\sqrt{2}} (|00\rangle + |11\rangle)$$

$$|\phi^-\rangle = \frac{1}{\sqrt{2}} (|00\rangle - |11\rangle) = (Z \otimes I) |\phi^+\rangle = (I \otimes Z) |\phi^+\rangle$$

$$|\psi^+\rangle = \frac{1}{\sqrt{2}} (|01\rangle + |10\rangle) = (X \otimes I) |\phi^+\rangle = (I \otimes X) |\phi^+\rangle$$

$$|\psi^-\rangle = \frac{1}{\sqrt{2}} (|01\rangle - |10\rangle) = (X \otimes I) |\phi^+\rangle = (I \otimes XZ) |\phi^+\rangle$$

We also write the four Bell states as

$$|\phi_{\alpha\beta}\rangle = (Z^\alpha X^\beta \otimes I) |\phi^+\rangle = (I \otimes X^\beta Z^\alpha) |\phi^+\rangle$$

$\alpha, \beta = 0, 1$

Chapter III: pg 22
Outcome probabilities for meas. outcome $|\phi_{\alpha\beta}\rangle$:

$$p_A = \text{tr}_B [|\phi\rangle\langle\phi|_{AB}] = \frac{1}{2} I_A \quad \leftarrow \text{state of A.}$$

$$p_{\alpha\beta} = \langle \phi_{\alpha\beta} | |XX\rangle_{A'} \otimes \frac{1}{2} I_A | \phi_{\alpha\beta} \rangle$$

$$= \frac{1}{2} \text{tr} [(|XX\rangle_{A'} \otimes I_A) | \phi_{\alpha\beta} \rangle \langle \phi_{\alpha\beta} |]$$

$$= \frac{1}{2} \text{tr}_{A'} [|XX\rangle_{A'} \cdot \underbrace{\text{tr}_A [| \phi_{\alpha\beta} \rangle \langle \phi_{\alpha\beta} |]}_{= \frac{1}{2} I_{A'}}]$$

$$= \frac{1}{2} \text{tr} [|XX\rangle_{A'} \cdot \frac{1}{2} I_{A'}]$$

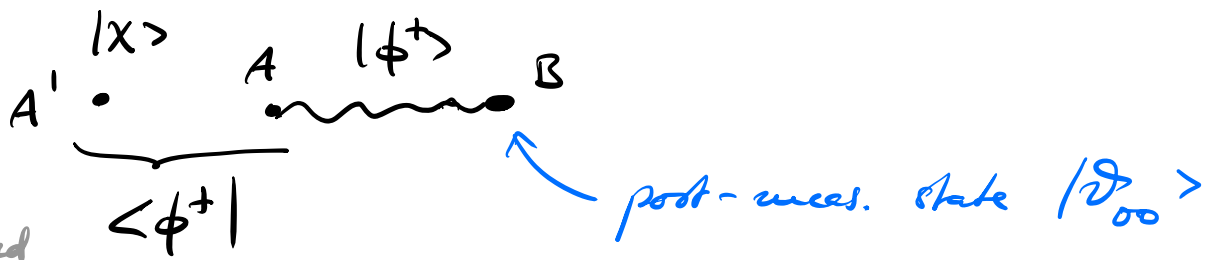
$$= \underline{\underline{\frac{1}{4}}}$$

$\Rightarrow$ equal probability $p_{\alpha\beta} = 1/4$ for all outcomes.

(This is good - if $p_{\alpha\beta}$ would depend on $|X\rangle$, it would reveal information on $|X\rangle$ and thus perturb the state!)

What is the state of B after the measurement?

i) Outcome $|\phi^+\rangle = |\phi_{00}\rangle$:



unnormalized

$$|\tilde{\mathcal{D}}_{00}\rangle = \langle \phi^+ |_{A'A} (|\chi\rangle_{A'} \otimes |\phi^+\rangle_{AB})$$

$$= \frac{1}{2} \left(\langle 00 |_{A'A} + \langle 11 |_{A'A} \right) \left((a|0\rangle_{A'} + b|1\rangle_{A'}) \otimes (|00\rangle_{AB} + |11\rangle_{AB}) \right)$$

$$= a \langle 0 |_A + b \langle 1 |_A$$

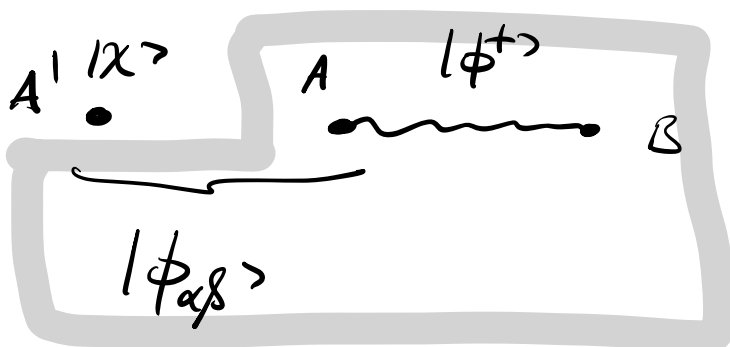
$$= \frac{1}{2} (a|0\rangle_B + b|1\rangle_B)$$

(*)

$\Rightarrow$ State $|\mathcal{D}_{00}\rangle = |\chi\rangle$ appears at B!

(works with 25% probability.)

ii) What about the other outcomes?



First consider $\langle \phi_{\alpha\beta} |_{A'A} | \phi^+ \rangle_{AB}$ — worked

gray above:

$$\begin{aligned}
 \langle \phi_{\alpha\beta} |_{A'A} | \phi^+ \rangle_{AB} &= \langle \phi^+ |_{A'A} (I_{A'} \otimes Z_A^\alpha X_A^\beta) | \phi^+ \rangle_{AB} \\
 &= \langle \phi^+ |_{A'A} (Z_A^\alpha X_A^\beta \otimes I_B) | \phi^+ \rangle_{AB} \\
 &= \langle \phi^+ |_{A'A} (I_A \otimes X_B^\beta Z_B^\alpha) | \phi^+ \rangle_{AB} \\
 &= X_B^\beta Z_B^\alpha \langle \phi^+ |_{A'A} | \phi^+ \rangle_{AB}
 \end{aligned}$$

Now combine with derivation $\otimes$ in part i)

$$\begin{aligned}
 |\tilde{\mathcal{Q}}_{\alpha\beta}\rangle &= \langle \phi_{\alpha\beta} |_{A'A} (| \chi \rangle_{A'} \otimes | \phi^+ \rangle_{AB}) \\
 &= X_B^\beta Z_B^\alpha \underbrace{\langle \phi^+ |_{A'A} (| \chi \rangle_{A'} \otimes | \phi^+ \rangle_{AB})}_{\stackrel{\otimes}{=} \frac{1}{2} | \chi \rangle_B} \\
 &= \underline{\underline{\frac{1}{2} X_B^\beta Z_B^\alpha | \chi \rangle_B}}.
 \end{aligned}$$

$\Rightarrow$ After A's measurement, B obtains $|\mathcal{D}_{\alpha\beta}\rangle = X^\beta Z^\alpha |\chi\rangle$ with probability $1/4$ each.

$\Rightarrow$ average state of B - without knowing meas. result - is $\frac{1}{4} \sum X^\beta Z^\alpha |\chi\rangle\langle\chi| Z^\alpha X^\beta = \frac{1}{2} I$.

i.e.: Bob has no information about $|\chi\rangle$
(in fact: same state as without meas.)

② A communicates meas. outcome (α, β) to B, and

③ B applies $(X^\beta Z^\alpha)^\dagger$ to their state

$\Rightarrow$ Bob obtains

$$(X^\beta Z^\alpha)^\dagger |\mathcal{D}_{\alpha\beta}\rangle = (X^\beta Z^\alpha)^\dagger (X^\beta Z^\alpha) |\chi\rangle = \underline{|\chi\rangle}$$

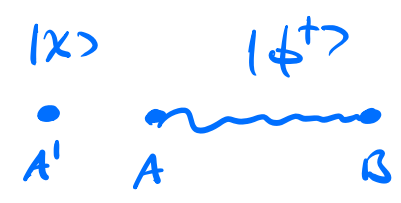
$\Rightarrow$ Bob obtains $|\chi\rangle$ with probability 1!

$\Rightarrow$ State $|\chi\rangle$ has been teleported to B.

Notes: • No faster-than-light communication
(avg. state of B is $\frac{1}{2} I$ prior to receiving (α, β) - which has finite transmiss. speed.)

- Communicating 1 qubit requires 1 "e-bit"
(= a max. entangled state $|\phi^+\rangle$ of 1+1 qubit)
+ 2 bits of classical communication ("c-bits")

Teleportation protocol - summary:



- ① Prepare A, A' in $|\phi_{\alpha\beta}\rangle$ state.
- ② Communicate (α, β) from A to B .
- ③ Apply $(X^\beta Z^\alpha)^\dagger$ on B .

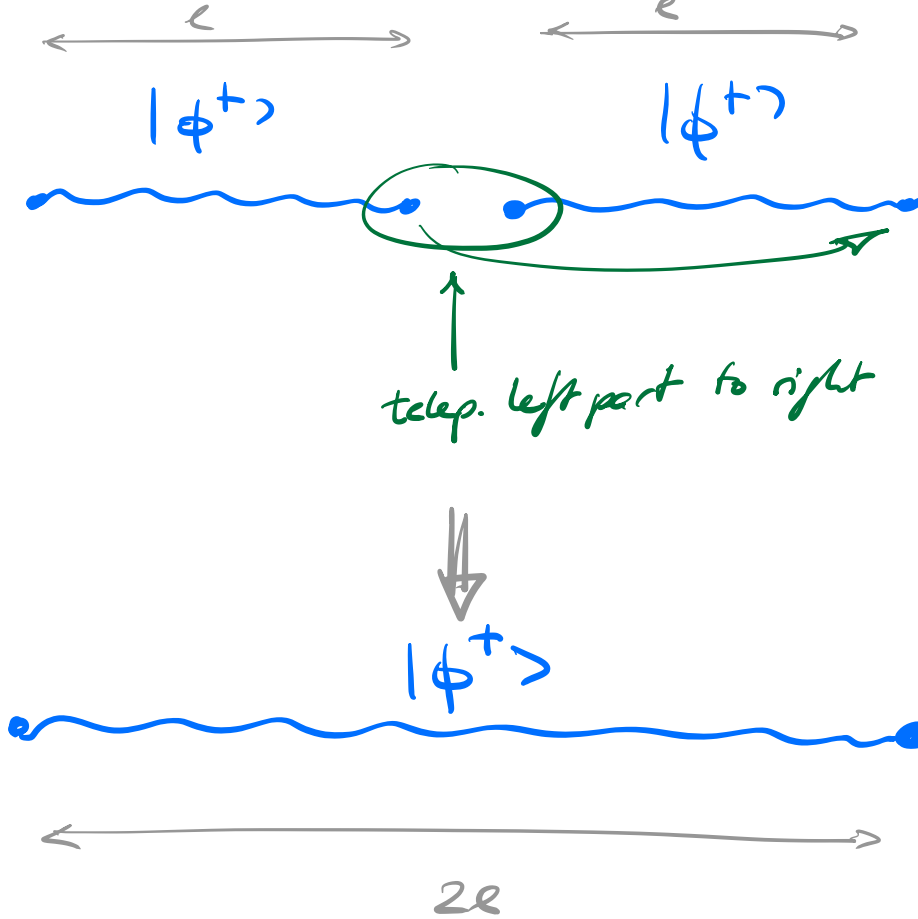
Can be straightforwardly generalized to $\mathbb{C}^d$.

One application of teleportation:

Quantum Repeaters

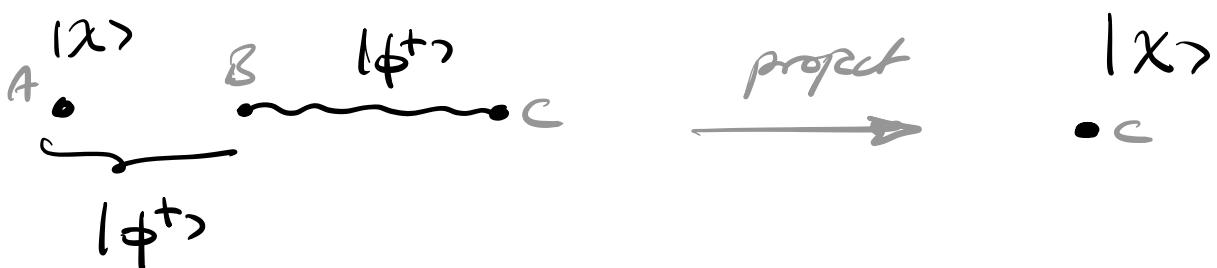
We can (reliably) create entanglement over distance $l \rightarrow$ can we create entanglement over distance $2l$?

(E.g.: Photon loss at const. rate $\rightarrow$ prob. to send half of an ent. pair over dist. l is $e^{-l/\xi}$.)



b) Relation between teleportation and the Choi-Jamiołkowski isomorphism

① Consider "postselected teleportation"



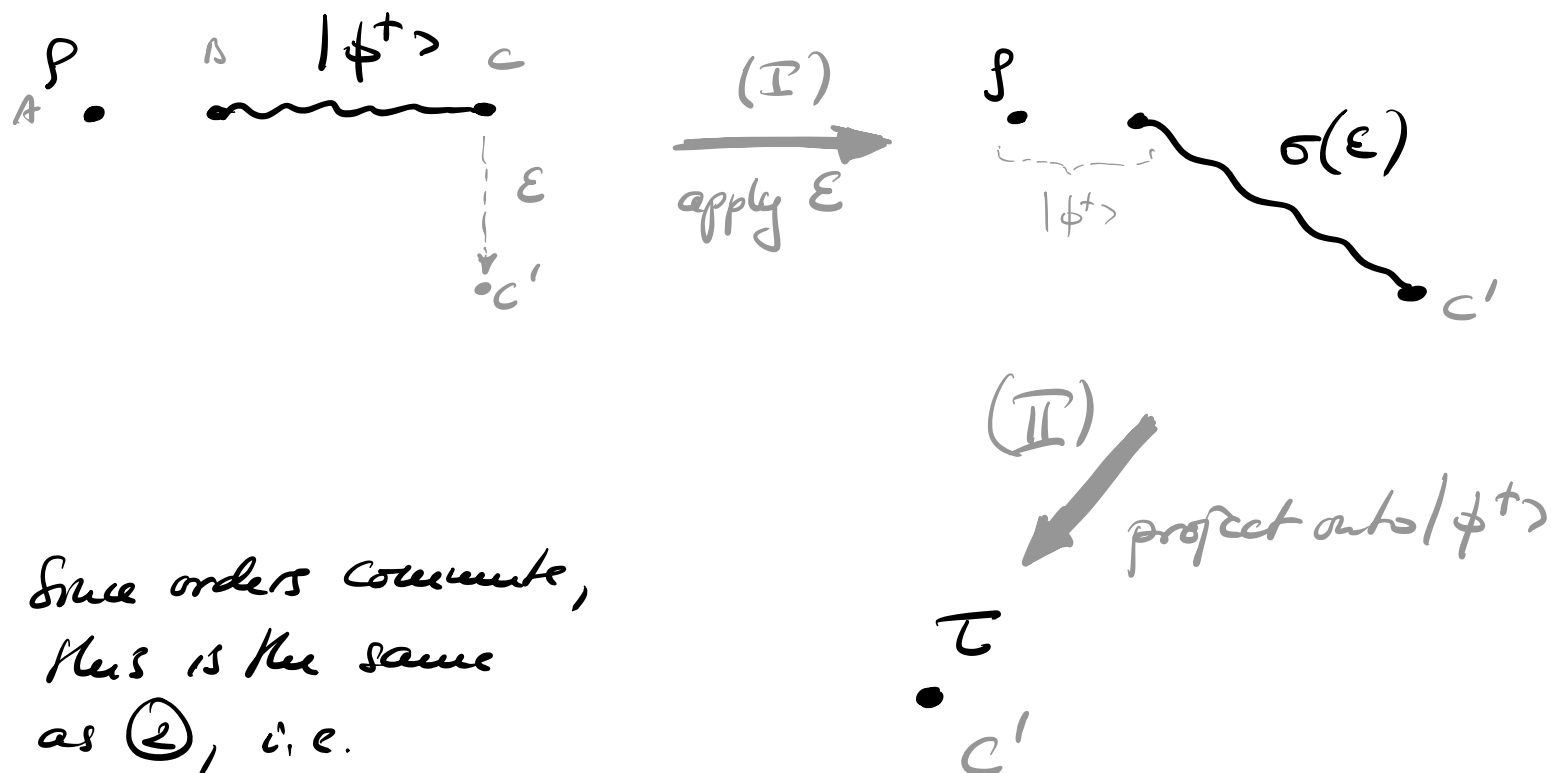
project onto $|\phi^+\rangle$: "postselected" measurement, i.e. we only consider this outcome

... so this is a complicated way of writing the identity map. Chapter III, pg 28

② Protocol for applying $p \mapsto E(p)$:



③ Now interchange the order of applying E and projecting — they commute (as they act on diff. systems), so this is the same map:



Since orders commute,
this is the same
as ②, i.e.

$$\tau = E(p)!$$

This is the Choi-Jamiołkowski isomorphism (!):

(I) is the $E \mapsto \sigma$ map

("apply E to half a max. entangled state")

(II) is the $\sigma \mapsto E$ map

("teleport σ through the Choi state")

c) Dense coding

Have seen:

- shared entanglement + class. channel $\rightarrow$ q. channel

$$1 \text{ ebit} + 2 \text{ cbit} \rightarrow 1 \text{ qubit}$$

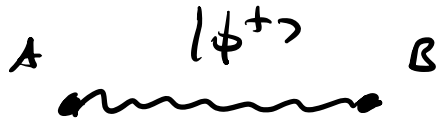
Can we do the converse? Use a quantum channel to transmit classical information?

Initially possible by encoding $0 \rightarrow |0\rangle$, $1 \rightarrow |1\rangle$

$$1 \text{ qubit} \rightarrow 1 \text{ cbit}$$

Chapter III, pg 30
Can we do better if we also share entanglement?

Dense coding (sometimes also "superdense coding"):



Idea: Encode two bits in $\{|\phi_{\alpha\beta}\rangle\}_{\alpha,\beta=0,1}$ (an ONB)

① A & B share $|\phi^+\rangle$.

② A can encode two bits α, β locally:

$$|\phi_{\alpha\beta}\rangle_{AB} = (Z_A^\alpha X_B^\beta \otimes I) |\phi^+\rangle_{AB}$$

i.e., A applies $Z^\alpha X^\beta$ to her part of $|\phi^+\rangle$.

③ A sends her part of the state to B via the quantum communication channel.

④ B measures in Bell basis $\{|\phi_{\alpha\beta}\rangle\}$ and recovers α and β .

shared ent. + q, channel $\rightarrow$ class. channel

1 cbit + 1 qubit $\rightarrow$ 2 cbit

d) Optimality of teleportation & dense coding

We can use the teleportation & dense coding protocol mutually to argue that both are optimal in terms of communication cost.

To this end, assume shared ent. is free (i.e.: this is not part of our cost function).

i) Assume we can teleport with $r < 2$ bits of class. communication per qubit sent (i.e., there are protocols to send k_q qubits w/ k_c class. bits s.t., $\frac{k_c}{k_q} \rightarrow r$).

Use this "hyper-teleportation" protocol to send quantum states in the (normal) dense coding prot.:

send $2n$ cbits

↓ dense coding

send n qubits

↓ "hyper-teleportation"

send rn cbits, $r < 2$

⇒ Can compress class. information (in the presence of entanglement).

⇒ Can iterate this to arbitrarily compress class. info
 - i.e., send n bits with $k \ll n$ bits - as long
 as we have free entanglement.

This is impossible! (intuitively, can also be formalized)

ii) Assume we can "hyper-dense-code" $2s > 2$ class.
 bits per qubit sent.

send $2s$ class. bits

↓

↓
send 1 qubit

↓ teleportation
send 2 bits

... and again, we can send 25 bits by only
transmitting 2 bits, etc. ...