

### 3. Applications of entanglement:

#### Teleportation and dense coding

##### a) Teleportation

Schup:



- A & B share entangled state  $|\phi^+\rangle_{AB} = \frac{1}{\sqrt{2}}(|00\rangle_{AB} + |11\rangle_{AB})$
- A has unknown quantum state  
 $|X\rangle_A = a|0\rangle_A + b|1\rangle_A$
- (Could e.g. also be part of a larger system  $\rightarrow$  memory!)
- A & B cannot (reliably\*) transmit quantum states, but can communicate classically "for free".

\* If the line is available, A & B can still use it to create entangled states  $|\phi^+\rangle$ , e.g. by repeat-until-success, or entanglement distillation ( $\rightarrow$  later!), or using "quantum repeaters" ( $\rightarrow$  later!)

Question: Can A get  $|X\rangle$  (safely) to B?

Problem: Any measurement of  $\langle X \rangle$  would <sup>Chapter III, pg 21</sup> reveal partial information, yet destroy state!

Solution: Quantum Teleportation!

Teleportation Protocol:

① A performs measurement on  $A'A$  in Bell Basis

$$|\phi^+\rangle = \frac{1}{\sqrt{2}} (|00\rangle + |11\rangle)$$

$$|\phi^-\rangle = \frac{1}{\sqrt{2}} (|00\rangle - |11\rangle) = (Z \otimes I) |\phi^+\rangle = (I \otimes Z) |\phi^+\rangle$$

$$|\psi^+\rangle = \frac{1}{\sqrt{2}} (|01\rangle + |10\rangle) = (X \otimes I) |\phi^+\rangle = (I \otimes X) |\phi^+\rangle$$

$$|\psi^-\rangle = \frac{1}{\sqrt{2}} (|01\rangle - |10\rangle) = (Z \otimes X) |\phi^+\rangle = (I \otimes Z \otimes X) |\phi^+\rangle$$

We also write the four Bell states as

$$|\phi_{\alpha\beta}\rangle = (Z^\alpha X^\beta \otimes I) |\phi^+\rangle = (I \otimes X^\beta Z^\alpha) |\phi^+\rangle$$

$\uparrow$   
 $\alpha, \beta = 0, 1$

Chapter III: pg 22

Outcome probabilities for meas. outcome  $|\phi_{\alpha\beta}\rangle$ :

$$P_A = \text{tr}_B \left[ |\phi^+ X \phi^+|_{AB} \right] = \frac{1}{2} I_A \quad \leftarrow \text{state of } A.$$

$$P_{\alpha\beta} = \langle \phi_{\alpha\beta} | |X X|_A \otimes \frac{1}{2} I_A | \phi_{\alpha\beta} \rangle$$

$$= \frac{1}{2} \text{tr} \left[ (|X X|_A \otimes I_A) | \phi_{\alpha\beta} X \phi_{\alpha\beta} | \right]$$

$$= \frac{1}{2} \text{tr}_{A'} \left[ |X X|_A \cdot \underbrace{\text{tr}_A \left[ | \phi_{\alpha\beta} X \phi_{\alpha\beta} | \right]}_{= \frac{1}{2} I_{A'}} \right]$$

$$= \frac{1}{2} \text{tr} \left[ |X X|_A \cdot \frac{1}{2} I_{A'} \right]$$

$$\underline{= \frac{1}{4}}$$

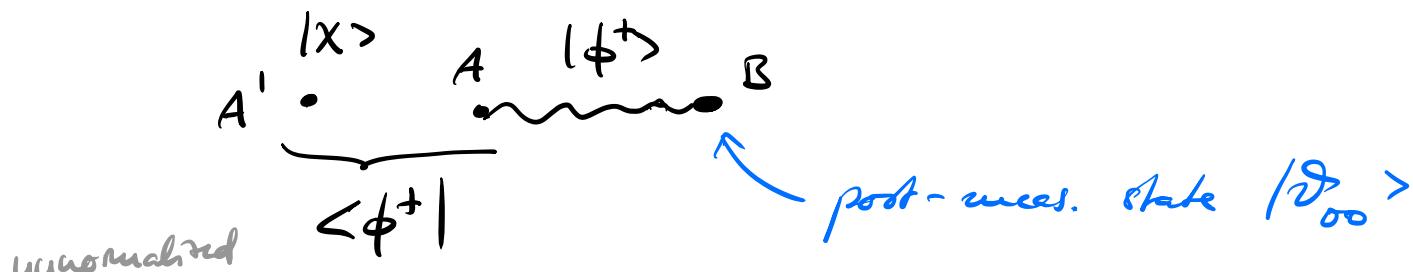
$\Rightarrow$  equal probability  $P_{\alpha\beta} = \frac{1}{4}$  for all outcomes.

(This is good - if  $P_{\alpha\beta}$  would depend on  $|X\rangle$ , it would reveal information on  $|X\rangle$  and thus perhaps the state!)

Chapter III, pg 23

What is the state of B after the measurement?

i) Outcome  $|\phi^+\rangle = |\phi_{00}\rangle$ :



$$|\tilde{D}_{00}\rangle = \langle\phi^+|_{A'A} (|X\rangle_{A'} \otimes |\phi^+\rangle_{AB})$$

$$= \frac{1}{2} (\underbrace{\langle 00|_{A'A} + \langle 11|_{A'A}}_{= a\langle 0|_A + b\langle 1|_A}) ((a|0\rangle_{A'} + b|1\rangle_{A'}) \otimes (|00\rangle_{AB} + |11\rangle_{AB}))$$

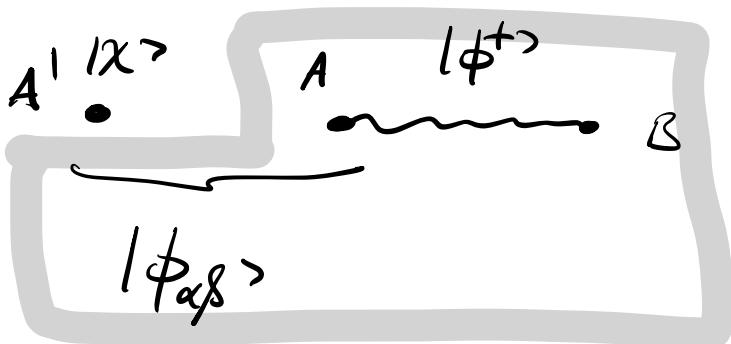
$$= \frac{1}{2} (a|0\rangle_B + b|1\rangle_B)$$

(~~✗~~)

⇒ State  $|\tilde{D}_{00}\rangle = |\chi\rangle$  appears at B!

(works with 25% probability.)

ii) What about the other outcomes?



First consider  $\langle \phi_{\alpha\beta} |_{A'A} |\phi^+ \rangle_{AB}$  - worked

gray above:

$$\begin{aligned}
 \langle \phi_{\alpha\beta} |_{A'A} |\phi^+ \rangle_{AB} &= \langle \phi^+ |_{A'A} (I_{A'} \otimes Z_A^\alpha X_A^\beta) |\phi^+ \rangle_{AB} \\
 &= \langle \phi^+ |_{A'A} (Z_A^\alpha X_A^\beta \otimes I_B) |\phi^+ \rangle_{AB} \\
 &= \langle \phi^+ |_{A'A} (I_A \otimes X_B^\beta Z_B^\alpha) |\phi^+ \rangle_{AB} \\
 &= X_B^\beta Z_B^\alpha \langle \phi^+ |_{A'A} |\phi^+ \rangle_{AB}
 \end{aligned}$$

Now combine with derivation  $\oplus$  in part i)

$$\begin{aligned}
 |\tilde{\phi}_{\alpha\beta}\rangle &= \langle \phi_{\alpha\beta} |_{A'A} (|x\rangle_{A'} \otimes |\phi^+ \rangle_{AB}) \\
 &= X_B^\beta Z_B^\alpha \underbrace{\langle \phi^+ |_{A'A} (|x\rangle_{A'} \otimes |\phi^+ \rangle_{AB})}_{\oplus} \\
 &= \frac{1}{2} X^\beta Z^\alpha |x\rangle_B
 \end{aligned}$$

⇒ After A's measurement, B obtains  $|D_{\alpha\beta}\rangle = X^{\beta} Z^{\alpha} |X\rangle$   
 Chapter 15, pg. 25

with probability  $\frac{1}{4}$  each.

⇒ average state of B — without knowing meas.  
 result is  $\frac{1}{4} \sum X^{\beta} Z^{\alpha} |X\rangle \langle X| Z^{\alpha} X^{\beta} = \frac{1}{2} I$ .

i.e.: Bob has no information about  $|X\rangle$

(in fact: same state as without meas.)

- ② A communicates meas. outcome  $(\alpha, \beta)$  to B, and
- ③ B applies  $(X^{\beta} Z^{\alpha})^+$  to their state

⇒ Bob obtains

$$(X^{\beta} Z^{\alpha})^+ |D_{\alpha\beta}\rangle = (X^{\beta} Z^{\alpha})^+ (X^{\beta} Z^{\alpha}) |X\rangle = \underline{|X\rangle}$$

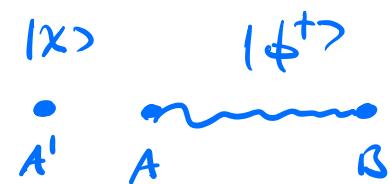
⇒ Bob obtains  $|X\rangle$  with probability 1!

⇒ State  $|X\rangle$  has been teleported to B.

Notes: • No faster-than-light communication  
 (avg. state of B is  $\frac{1}{2} I$  prior to receiving  $(\alpha, \beta)$  — which has finite transm. speed.)

- Communicating 1 qubit requires <sup>Chapter "Entanglement"</sup>  $2^{2^k}$  bits (= a max. entangled state  $|\phi^+\rangle$  of 1+1 qubit) + 2 bits of classical communication ("c-bits")

## Teleportation protocol - summary:



- ① Receive  $A, A'$  in  $|\phi_{\alpha\beta}\rangle$  basis.
- ② Communicate  $(\alpha, \beta)$  from  $A$  to  $B$ .
- ③ Apply  $(X^{\beta} Z^{\alpha})^+$  on  $B$ .

Can be straightforwardly generalized to  $\mathbb{C}^d$ .

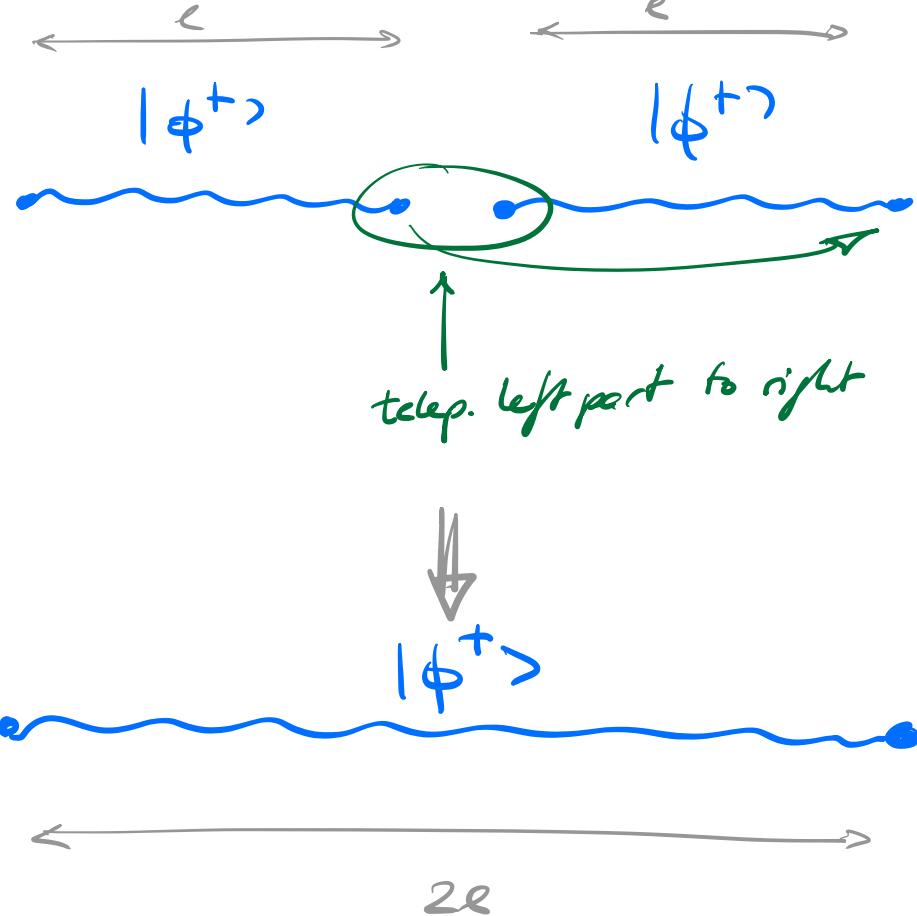
One application of teleportation:

## Quantum Repeaters

We can (reliably) create entanglement over distance  $\ell \rightarrow$  can we create entanglement over distance  $2\ell$ ?

(E.g.: Photon loss at const. rate  $\rightarrow$  prob. to

send half of an ent. pair over dist.  $\ell$  is  $e^{-\epsilon/\xi}$ .)



### b) Relation between teleportation

and the Choi-Jamiołkowski isomorphism

- ① Consider "postselected teleportation"



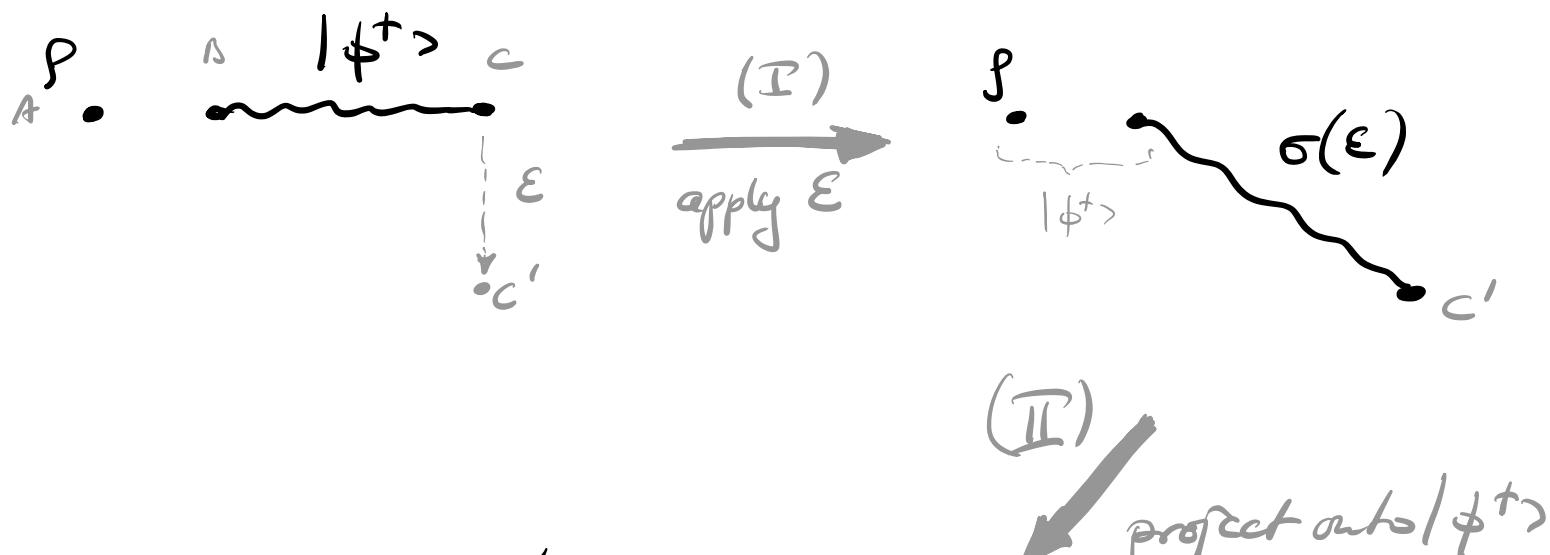
project onto  $|\phi^+\rangle$ : "postselected" measurement, i.e. we only consider this outcome

... so this is a complicated way of writing the identity map. Chapter III, pg 28

② Protocol for applying  $\rho \mapsto E(\rho)$ :



③ Now interchange the order of applying  $E$  and projecting — they commute (as they act on diff. systems), so this is the same res:



Since orders commute,  
this is the same  
as ②, i.e.

$$\tau = E(\rho)!$$

$\tau$   
 $c'$

This is the Choi-Jamiołkowski isomorphism (?):

(I) is the  $\mathcal{E} \mapsto \sigma$  map

("apply  $\mathcal{E}$  to half a max. entangled state")

(II) is the  $\sigma \mapsto \mathcal{E}$  map

("teleport  $\sigma$  through the Choi state")

### c) Dense coding

Have seen:

- shared entanglement + class. channel  $\rightarrow$  q. channel

$$1\text{ebit} + 2\text{ cbits} \rightarrow 1\text{qubit}$$

Can we do the converse? Use a quantum channel to transmit classical information?

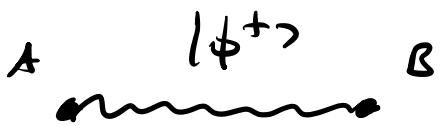
Trivially possible by encoding  $0 \rightarrow |0\rangle$ ,  $1 \rightarrow |1\rangle$

$$1\text{qubit} \rightarrow 1\text{ebit}$$

Can we do better if we also share certain ~~secret~~?

Chapter III, pg 30

Dense coding (sometimes also "hypercold coding"):



Idea: Encode two bits in  $\{|\phi_{\alpha\beta}\rangle\}_{\alpha,\beta=0,1}$  (an ONB)

① A & B share  $|\phi^+\rangle$ .

② A can encode two bits  $\alpha, \beta$  locally:

$$|\phi_{\alpha\beta}\rangle_{AB} = (Z_A^\alpha X_B^\beta \otimes I) |\phi^+\rangle_{AB}$$

i.e., A applies  $Z^\alpha X^\beta$  to her part of  $|\phi^+\rangle$ .

③ A sends her part of the state to B via the quantum communication channel.

④ B measures in Bell basis  $\{|\phi_{\alpha\beta}\rangle\}$  and recovers  $\alpha$  and  $\beta$ .

shared ent. + q. channel  $\rightarrow$  class. channel

1 ebit + 1 qubit  $\rightarrow$  2 cbit

### d) Optimality of teleportation & dense coding

We can use the teleportation & dense coding protocol mutually to argue that both are optimal in terms of communication cost.

To this end, assume shared ent. is free (i.e.: this is not part of our cost function).

- i) Assume we can teleport with  $r < 2$  bits of class. communication per qubit sent (i.e., there are protocols to send  $k_q$  qubits w/  $k_c$  class. bits s.t.  $\frac{k_c}{k_q} \rightarrow r$ ).

Use this "hyper-teleportation" protocol to send quantum states in the (normal) dense coding prot.:

Send 2n cbits

↓  
dense coding

Send n qubits

↓  
"hyper-telportation"

Send r n cbits,  $r < 2$

⇒ Can compress class. information (in the presence of entanglement).

⇒ Can iterate this to arbitrarily compress class. info  
- i.e., send n bits with  $k \ll n$  bits - as long  
as we have free entanglement.

This is impossible! (Intuitively, can also be formalized.)

ii) Assume we can "hyper-dense-code"  $2s \geq 2$  class.  
bits per qubit sent.

Send  $2s$  class. bits

↓  
send 1 qubit

↓  
teleportation  
send 2 cbs

... and again, we can send 2s bits by only  
transmitting 2 bits, etc. ...