

Quantum teleportation & dense coding

Assume A has an (unknown) state $|x\rangle \in \mathbb{C}^2$.

Her goal is to transfer this state to B,

but only with classical communication.

If she had many copies of $|x\rangle$, she could

measure many times, gather statistics, and

thus obtain the exact form $|x\rangle = \chi_0|0\rangle + \chi_1|1\rangle$,

(also known as tomography)

then communicate χ_0, χ_1 to Bob, who can then prepare the state.

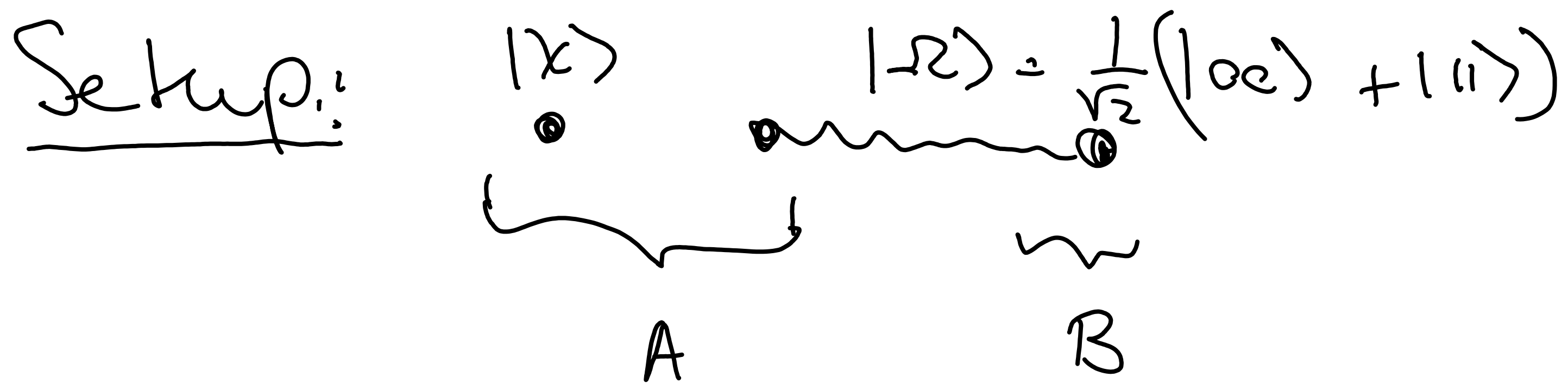
What if she has only a single copy of $|x\rangle$?

Problem: measurement destroys the state,

she can't obtain classical description.

However, if A and B share an

entangled state, then A can still transfer $|x\rangle$ to B. Only 2 bits are needed!



To notice: $\mathbb{C}^2 \otimes \mathbb{C}^2$ has the following

basis:

$$|R_{00}\rangle = |R\rangle = \frac{1}{\sqrt{2}}(|00\rangle + |11\rangle)$$

$$|R_{01}\rangle = (1 \otimes Z)|R\rangle = \frac{1}{\sqrt{2}}(|00\rangle - |11\rangle)$$

$$|R_{10}\rangle = (1 \otimes X)|R\rangle = \frac{1}{\sqrt{2}}(|01\rangle + |10\rangle)$$

$$|R_{11}\rangle = (1 \otimes XZ)|R\rangle = \frac{1}{\sqrt{2}}(|01\rangle - |10\rangle)$$

$$\left. \begin{array}{l} |R_{\alpha\beta}\rangle = X^{\alpha} Z^{\beta} \cdot |R\rangle \end{array} \right\}$$

We can thus write $|x\rangle \otimes |R\rangle \in \mathbb{C}^2 \otimes \mathbb{C}^2 \otimes \mathbb{C}^2$

$$|x\rangle \otimes |R\rangle = \sum_{\alpha\beta} |R_{\alpha\beta}\rangle \otimes |\varphi_{\alpha\beta}\rangle, \text{ for}$$

some $|\varphi_{\alpha\beta}\rangle$. How to calculate $|\varphi_{\alpha\beta}\rangle$?

(a) Brute force

(b) Notice that

$$|\psi_{\alpha\beta}\rangle = (\langle \Omega_{\alpha\beta} | \otimes \text{id}) (|\chi\rangle \otimes |\Omega\rangle)$$

$$= (\langle \Omega | \otimes \text{id}) (\text{id} \otimes Z^\beta X^\alpha \otimes \text{id}) (|\chi\rangle \otimes |\Omega\rangle)$$

$$= (\langle \Omega | \otimes \text{id}) (\text{id} \otimes \text{id} \otimes X^\alpha Z^\beta) (|\chi\rangle \otimes |\Omega\rangle)$$

$$\uparrow$$

$$(X \otimes \text{id}) |\Omega\rangle = (\text{id} \otimes X) |\Omega\rangle$$

$$(Z \otimes \text{id}) |\Omega\rangle = (\text{id} \otimes Z) |\Omega\rangle$$

$$= X^\alpha Z^\beta \cdot \underbrace{(\langle \Omega | \otimes \text{id}) (\text{id} \otimes |\Omega\rangle)}_{|\chi\rangle}$$

⊛

$$(\langle \Omega | \otimes \text{id}) (\text{id} \otimes |\Omega\rangle) = \frac{1}{2} (\langle 00 | \otimes \text{id} + \langle 11 | \otimes \text{id}).$$

$$\cdot (\text{id} \otimes |00\rangle + \text{id} \otimes |11\rangle) = \frac{1}{2} (|0\rangle\langle 0| + |1\rangle\langle 1|),$$

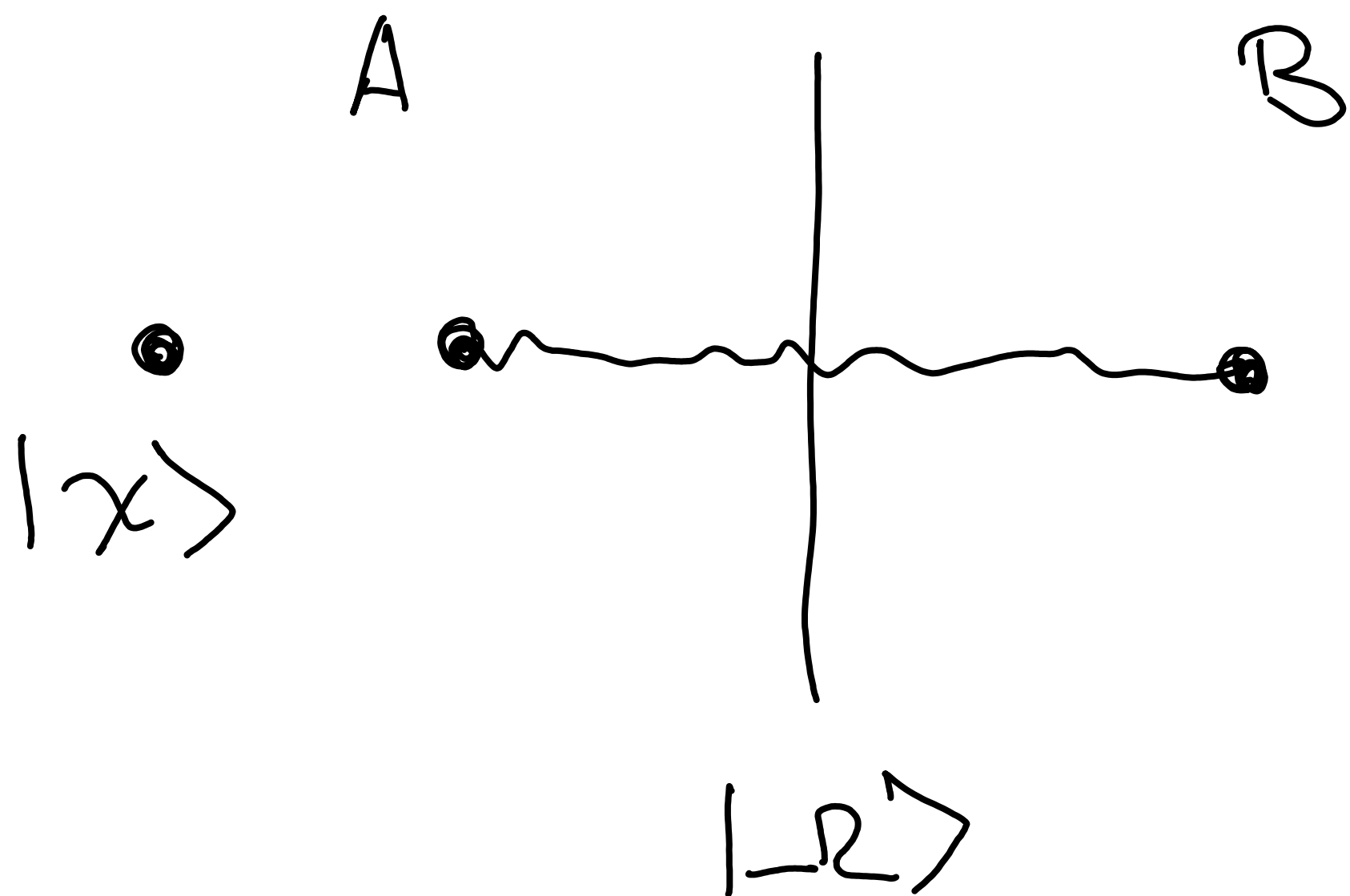
So this is just the identity!

We thus get that

$$|\chi\rangle \otimes |\Omega\rangle = \sum_{\alpha\beta} |\Omega_{\alpha\beta}\rangle \otimes X^\alpha Z^\beta |\chi\rangle.$$

Missing: ⊛

Teleportation:



3-partite system : $\mathcal{H}_{A_1} \otimes \mathcal{H}_{A_2} \otimes \mathcal{H}_B$,

$$|\Psi_{ABC}\rangle = |\chi\rangle \otimes |R\rangle$$

$$\mathcal{H}_{A_1} \cong \mathbb{C}^2$$

$$\mathcal{H}_{A_2} \otimes \mathcal{H}_B \cong \mathbb{C}^2 \otimes \mathbb{C}^2$$

Goal : transfer $|\chi\rangle$ to B's side.

Rules : A can manipulate her side (A_1 & A_2), B his side (B), but they can't act globally. However, they are allowed to communicate.

Can they achieve their goal? Yes!

For that, let us rewrite the state w.r.t. the $A_1 A_2 | B$ bipartition. Actually, we want a specific ONB on $A_1 A_2$:

$$|-\mathcal{R}_{\alpha\beta}\rangle = (1 \otimes X^\alpha Z^\beta) |-\mathcal{R}\rangle$$

This is called the Bell basis.

$$|-\mathcal{R}_{00}\rangle = \frac{1}{\sqrt{2}} (|00\rangle + |11\rangle)$$

$$|\mathcal{R}_{01}\rangle = \frac{1}{\sqrt{2}} (|00\rangle - |11\rangle)$$

$$|\mathcal{R}_{10}\rangle = \frac{1}{\sqrt{2}} (|01\rangle + |10\rangle)$$

$$|-\mathcal{R}_{11}\rangle = \frac{1}{\sqrt{2}} (|01\rangle - |10\rangle)$$

So we want $(\chi_{\alpha\beta})$ s.t.

$$|\Psi_{A_1 A_2 B}\rangle = |\chi\rangle \otimes |-\mathcal{R}\rangle = \sum_{\alpha\beta} |-\mathcal{R}_{\alpha\beta}\rangle \otimes |\chi_{\alpha\beta}\rangle.$$

As $|-\mathcal{R}_{\alpha\beta}\rangle$ form an ONB,

$$|\chi_{\alpha\beta}\rangle = \underbrace{\langle -\mathcal{R}_{\alpha\beta} | \otimes \text{id} }_{\substack{\mathcal{H}_{A_1} \otimes \mathcal{H}_{A_2} \rightarrow \mathbb{C} \\ \mathcal{H}_{A_1} \otimes \mathcal{H}_{A_2} \otimes \mathcal{H}_B \rightarrow \mathcal{H}_B}} \underbrace{|\Psi_{A_1 A_2 B}\rangle}_{\substack{\mathcal{H}_{A_1} \otimes \mathcal{H}_{A_2} \otimes \mathcal{H}_B}}$$

(Also could be noted as $\langle -\mathcal{R}_{\alpha\beta} |_{A_1 A_2} | \Psi_{A_1 A_2 B} \rangle$)

— here the index A, A_2 means that the lin. fun.
 $\langle -\mathcal{R}_{\alpha\beta} | : \mathbb{C}^2 \otimes \mathbb{C}^2 \rightarrow \mathbb{C}$ acts on the A, A_2 component)

So

$$|\chi_{\alpha\beta}\rangle = (\langle -\mathcal{R} | \otimes \mathbb{1}) (\mathbb{1} \otimes Z^\beta X^\alpha \otimes \mathbb{1}) (|\chi\rangle \otimes |-\mathcal{R}\rangle)$$

$$\langle -\mathcal{R}_{\alpha\beta} | = (|-\mathcal{R}_{\alpha\beta}\rangle)^\dagger = \left((\mathbb{1} \otimes X^\alpha Z^\beta) |-\mathcal{R}\rangle \right)^\dagger$$

$$\langle -\mathcal{R}_{\alpha\beta} | \otimes \mathbb{1}$$

Using now $(\mathbb{0} \otimes \mathbb{1}) |-\mathcal{R}\rangle = (\mathbb{1} \otimes \mathbb{0}^T) |-\mathcal{R}\rangle$,
 with $\mathbb{0} = Z^\beta X^\alpha$, we obtain

$$\begin{aligned} |\chi_{\alpha\beta}\rangle &= (\langle -\mathcal{R} | \otimes X^\alpha Z^\beta) (|\chi\rangle \otimes |-\mathcal{R}\rangle) \\ &= X^\alpha Z^\beta \cdot |\chi_{00}\rangle \end{aligned}$$

Here

$$|\chi_{00}\rangle = (\langle -\mathcal{R} | \otimes \mathbb{1}) (|\chi\rangle \otimes |-\mathcal{R}\rangle)$$

$$\begin{aligned} &= \frac{1}{2} \sum_{ij} (\langle i | \otimes \langle i | \otimes \mathbb{1}) (|\chi\rangle \otimes |j\rangle \otimes |j\rangle) \\ &= \frac{1}{2} \sum_{ij} \langle i | \chi \rangle \underbrace{\langle i | j \rangle}_{\delta_{ij}} \cdot |j\rangle = \frac{1}{2} |\chi\rangle. \end{aligned}$$

We can thus write

$$|\chi\rangle \otimes |\Omega\rangle = \frac{1}{2} \sum_{\alpha, \beta} |\Omega_{\alpha\beta}\rangle \otimes X^\alpha Z^\beta |\chi\rangle.$$

So we can do the following to transfer $|\chi\rangle$ to Bob's side:

- 1) A measures her 2 qubits in the $|\Omega_{\alpha\beta}\rangle$ basis (i.e., w/ meas. operators $\{|\Omega_{\alpha\beta}\rangle\langle\Omega_{\alpha\beta}| \}_{\alpha, \beta \in \{0, 1\}}$).

For outcome $\alpha\beta$, the post-meas. state is $|\Omega_{\alpha\beta}\rangle \otimes X^\alpha Z^\beta |\chi\rangle$.

- 2) She communicates the outcome α and β to Bob.

- 3) Bob knows that after measurement he has the state $X^\alpha Z^\beta |\chi\rangle$, so applies $Z^\beta X^\alpha$ to get back $|\chi\rangle$.

We can further break down

"A measures in the $|\Omega_{\alpha\beta}\rangle$ basis":

① Find U st. $U|\alpha\beta\rangle = |\Omega_{\alpha\beta}\rangle$

② Instead of measuring w/ $|\Omega_{\alpha\beta}\rangle$, first apply U^\dagger , then measure $|\alpha\beta\rangle$, i.e. comp. basis in both qubits, then rotate back w/ U . We don't really need A 's part of the state, so we can skip this part.

Note: if $Q_i = U P_i U^\dagger$, then

$$* \text{tr}\{Q_i \rho\} = \text{tr}\{P_i U^\dagger \rho U\}$$

$$* Q_i \rho Q_i^\dagger = U P_i U^\dagger \rho U P_i^\dagger U^\dagger.$$

Finally note that

$$|\Omega_{\alpha\beta}\rangle = \text{CNOT}(H \otimes I) |\alpha\beta\rangle :$$

$$\uparrow \quad \uparrow$$
$$\frac{1}{\sqrt{2}} \begin{pmatrix} 1 & 1 \\ 1 & -1 \end{pmatrix}, \quad HX = 2H$$

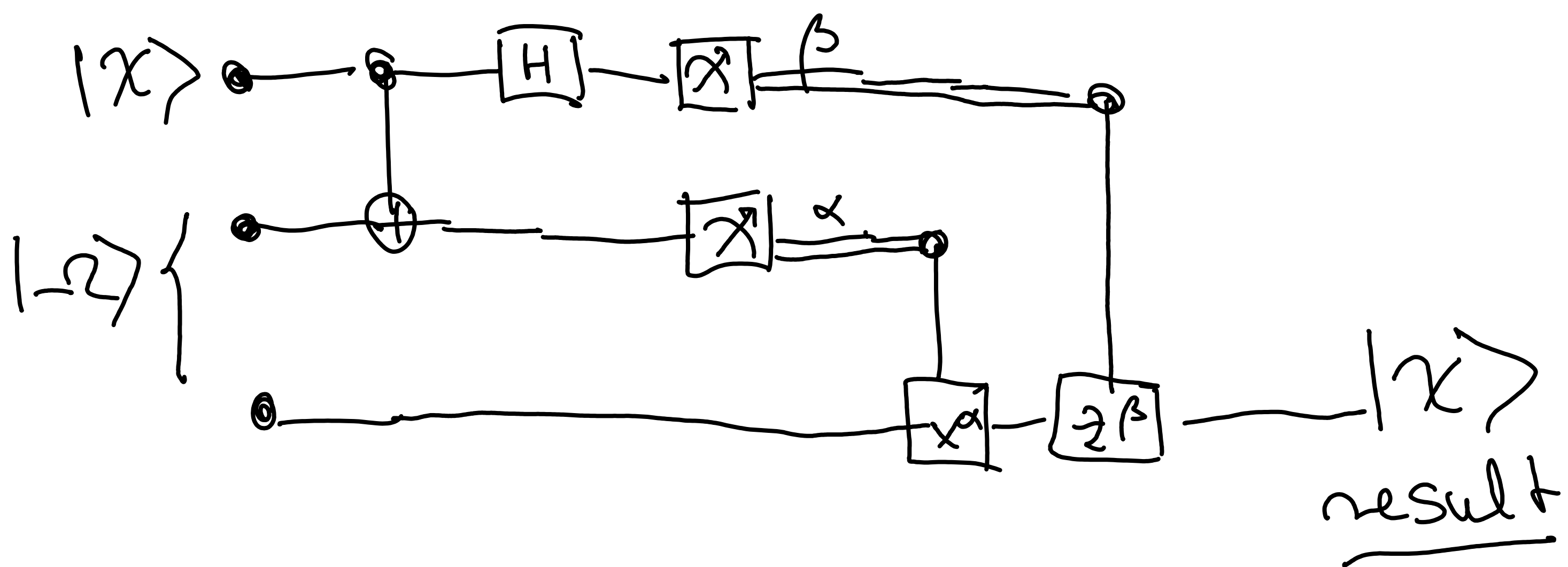
$$|0\rangle\langle 0| \otimes I + |1\rangle\langle 1| \otimes X$$

$$\begin{aligned}
\text{CNOT}(H \otimes I) |\beta\alpha\rangle &= \text{CNOT}(H \otimes I) (X^\beta \otimes X^\alpha) |00\rangle \\
&= \text{CNOT}(Z^\beta \otimes X^\alpha) (H \otimes I) |00\rangle \\
&= (Z^\beta \otimes X^\alpha) \underbrace{\text{CNOT}(H \otimes I) |00\rangle}_{\frac{1}{\sqrt{2}}(|00\rangle + |10\rangle)} \\
&= (I \otimes X^\alpha Z^\beta) |-\rangle.
\end{aligned}$$

$\frac{1}{\sqrt{2}} |00\rangle + |11\rangle = |-\rangle$

So if we first undo CNOT, then $(H \otimes I)$, we can measure in the comp. basis!

Final protocol:



Notice: We did Local Operations
assisted with Classical Communication,
LOCC.

Such operations should not increase
entanglement, as they are essentially classical.
In our case, we have started from

$$|\chi\rangle \otimes |\Omega\rangle \in (\mathcal{H}_{A_1} \otimes \mathcal{H}_{A_2}) \otimes \mathcal{H}_B,$$

which is entangled (wrt. A-B cut!).

After the operation, if we measure
 $\alpha\beta$, we end up in $|\alpha\beta\rangle \otimes |\chi\rangle$.

Actually, we can create $U_{\alpha\beta} \in \mathcal{H}_{A_1} \otimes \mathcal{H}_{A_2}$. St.

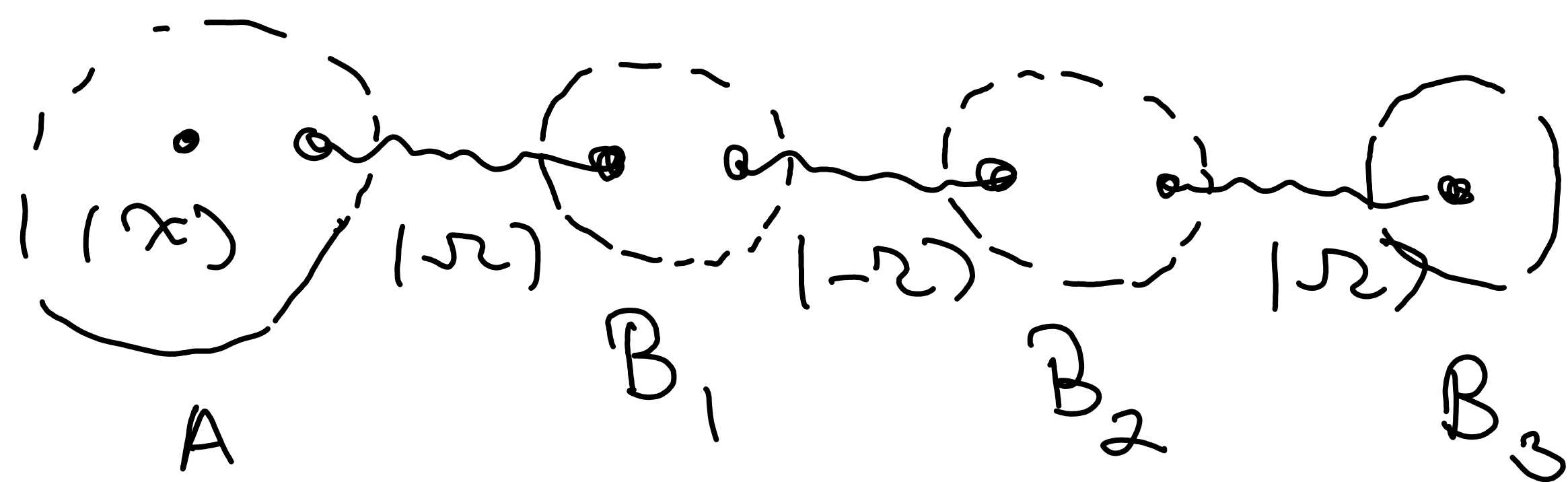
we always end up in the same state,

e.g. $|00\rangle \otimes |\chi\rangle$.

This state is NOT entangled wrt A/B cut,

LOCC can destroy entanglement.

Remark: teleportation can be used to transfer a state to large distances



We can transfer $|\chi\rangle$ to B₃ by repeated teleportation! This is called quantum repeater.

Remark II: Observe how similar this is to the Choi-Jamiołkowski isomorphism (when we get the CE outcome): there also we "teleport ρ " through the max. ent. state.

Dense coding

If A & B share an entangled state, then A can send B two classical bits of info via sending a single qubit.

1] A & B share $|\Omega\rangle = |\Phi^+\rangle = \frac{1}{\sqrt{2}}(|00\rangle + |11\rangle)$.

2] A prepares $|\Omega_{\alpha\beta}\rangle = (Z^\beta X^\alpha \otimes I)|\Omega\rangle$

3] After sending A's part to Bob, Bob can measure $|\Omega_{\alpha\beta}\rangle$, so by sending 1 qubit; A could send 2 classical bits information.

Remark: Needed a pre-shared max. ent. state, so all together 2 qubit transfer is needed. But 1 qubit can be transferred before the message.